

Distributed On-demand MAC Scheduling for Underwater Acoustic Networks

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Presented by Zheng Peng

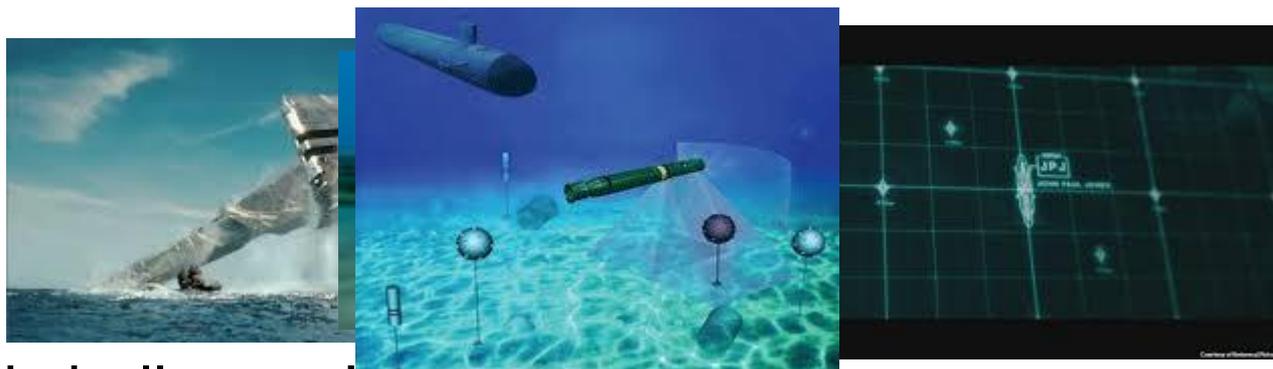
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Outline

- Background
- Motivation: Heavy Collisions
- Implication for Practical MAC Design
- Proposed Solution: Distributed On-demand Scheduling
- Simulation Results
- Conclusion

Underwater Acoustic Networks (UANs)

- Wide range of applications!



- Grand challenges!

- Acoustic communication

- Slow propagation speed

- sound speed in water: ~ 1500m/s vs. radio speed: 2×10^8

- Low available bandwidth

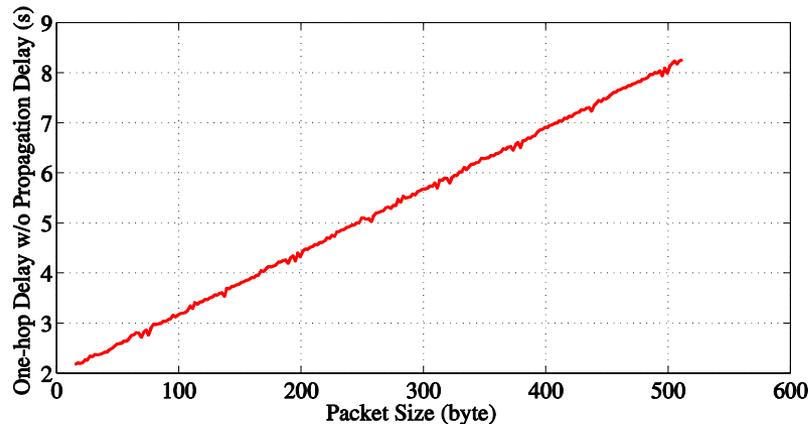
- acoustic: several kbps vs. radio: tens or hundreds of Mbps

- Dynamic environment

- Water current ...

Motivation – Acoustic Modem Characteristics

➤ One-hop delay with Benthos Modem



$$D = \frac{L}{667} + 1.9615$$

D : one-hop delay without propagation delay, in second
 L : The size of packet, in bits

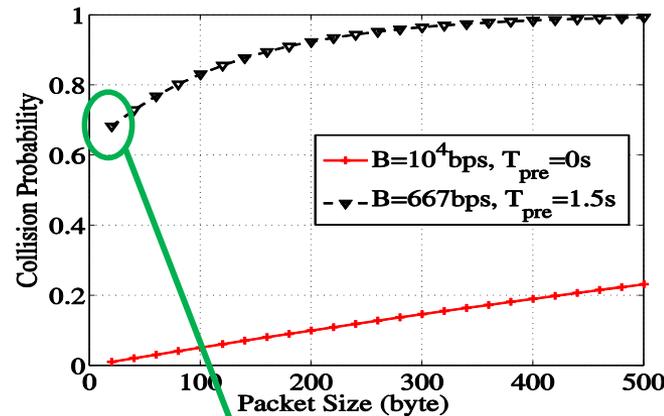
➤ Observed characteristics

➤ **Extremely low** effective transmission rate (conventionally 10 kbps)

➤ **Long** preamble : 1.5 s (conventionally ignored)

Impact on collision probability

- Collision probability P_c of ALOHA with varying packet size:



- Observation

- Significantly heavier collisions when considering the real modem characteristics

- Even short packet of 20 bytes collide with probability 0.68

Are control packets still effective for collision avoidance?

Implication for practical MAC design

- Contention is costly!
 - Random access-based MAC
 - Long transmission time cause severe collisions
 - Reservation-based MAC
 - Principle:
 - Control packet is very short and thus effective for collision avoidance
 - Radio Networks: works well
 - Underwater Acoustic Networks:
 - Control packet is very long
- 
- Key feature of practical MAC design:
COLLISION FREE, even for **control** packets

Distributed On-demand Scheduling (DOS)

- Both **control** and **data** packets are collision-free
 - Guarantees a high performance (contentions are costly)
- Distributed
 - Expensive to collect required information for a centralized scheduling algorithm
- On-demand
 - Nodes' traffic requirements are dynamic
- Pure scheduling only
 - Not modem-specific:
 - No CDMA
 - No power adjustment
 - No multi-user receiver system or multi-channel

Basic Ideas

- Cluster based network structure
 - Group the nodes into clusters, using existing clustering algorithms

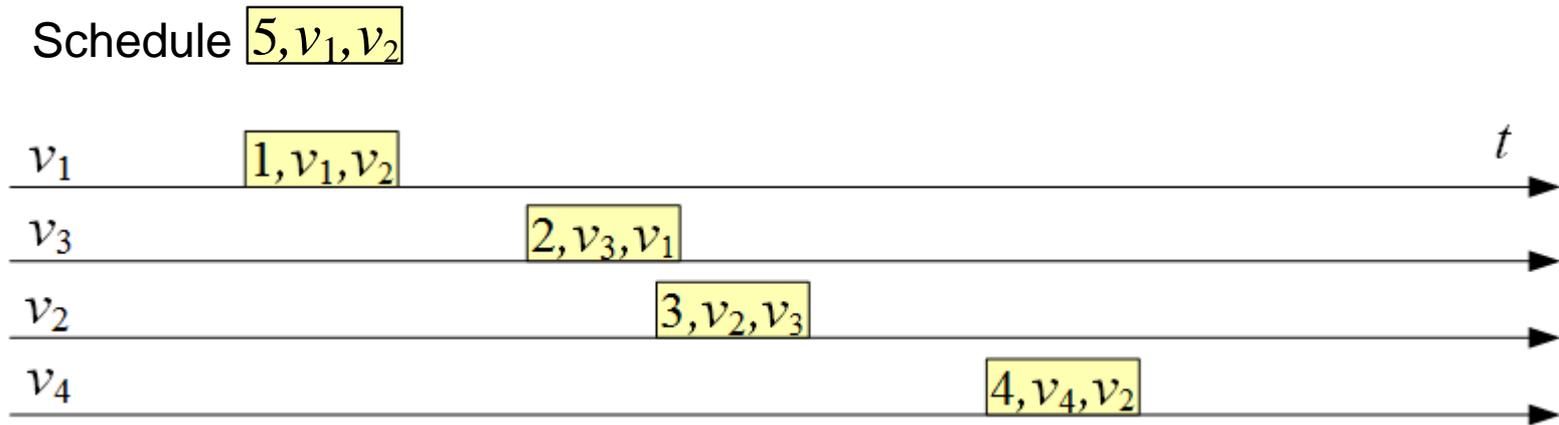
- The data transmissions within a cluster are scheduled by cluster heads
 - Cluster members contact the cluster head to request data transmissions

- Schedule exchanges between neighboring cluster heads
 - Every cluster head aware of what's going on in the neighboring cluster

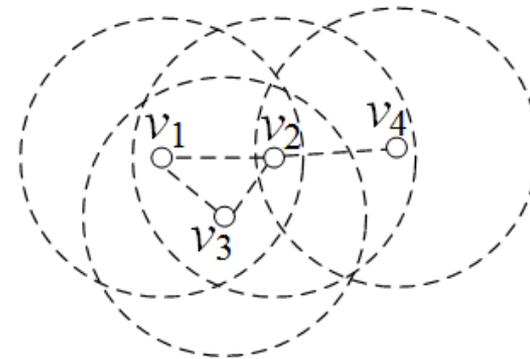
- Two hop neighbors' up-to-date tx schedule is enough for a node's new collision-free schedule

How to schedule a packet transmission?

➤ Example:

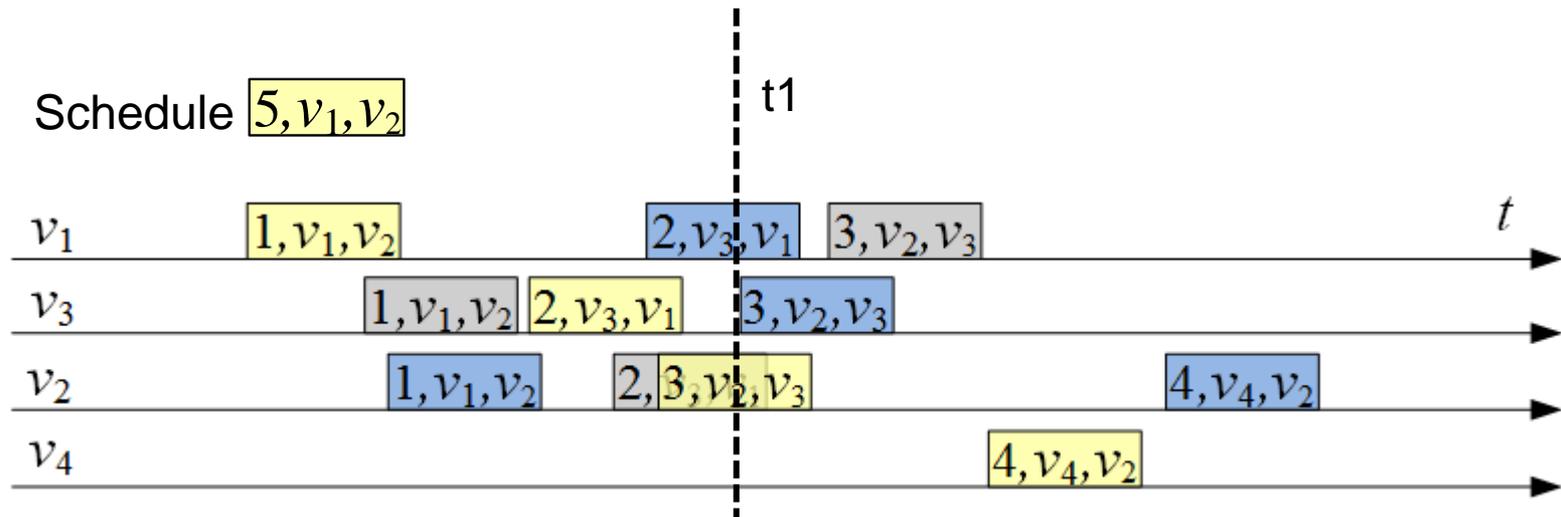


- $\boxed{1, v_1, v_2}$ Event of transmitting a packet from v_1 to v_2
- $\boxed{1, v_1, v_2}$ Event of receiving a packet from v_1 to v_2
- $\boxed{1, v_1, v_2}$ Event of overhearing a packet from v_1 to v_2
- $\boxed{}$ Newly scheduled TX slot

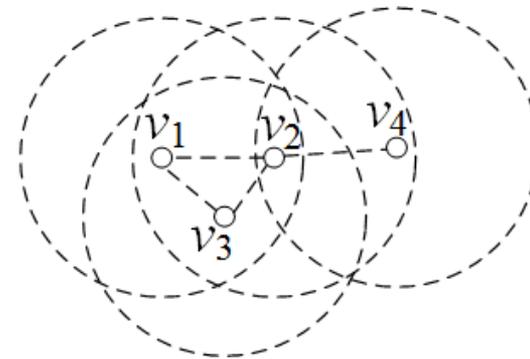


How to schedule a packet transmission?

➤ Corresponding events and sequence

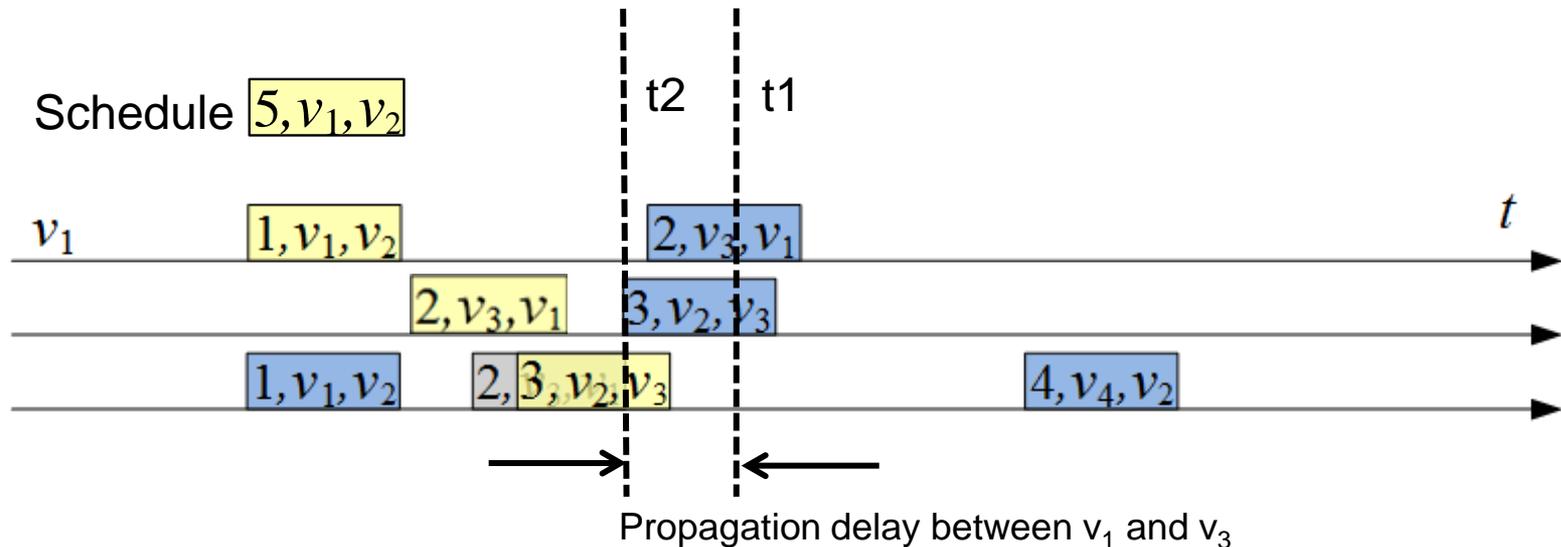


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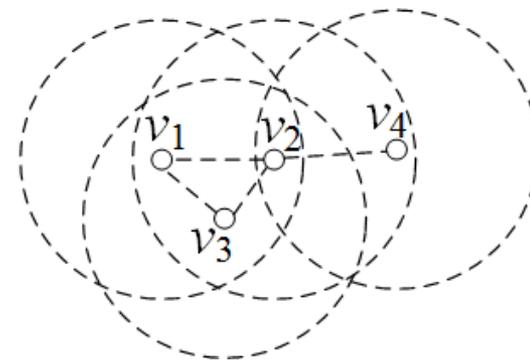


How to schedule a packet transmission?

➤ Time alignment and mapping



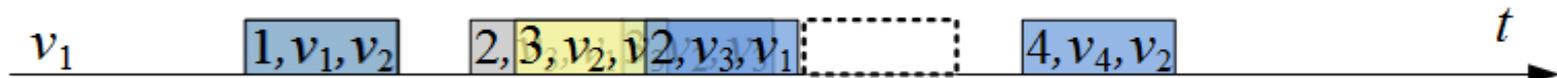
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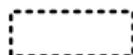


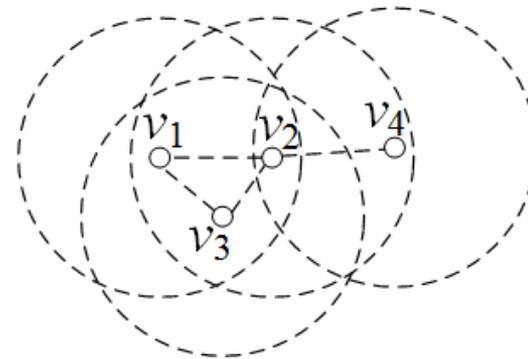
How to schedule a packet transmission?

- Combine all events to v_1 's timeline and find a time slot for packet #5

Schedule $5, v_1, v_2$

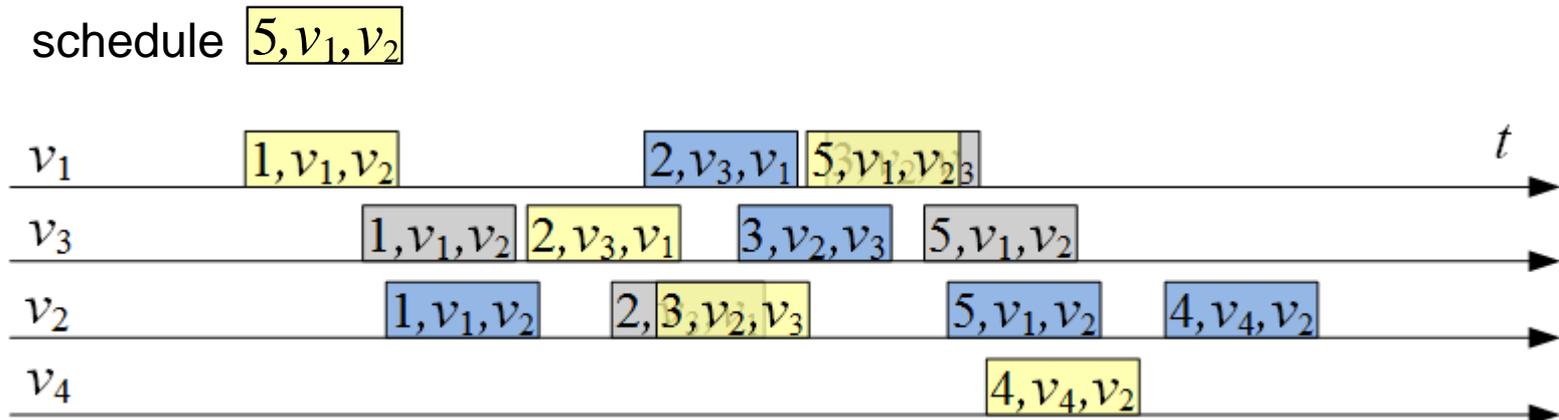


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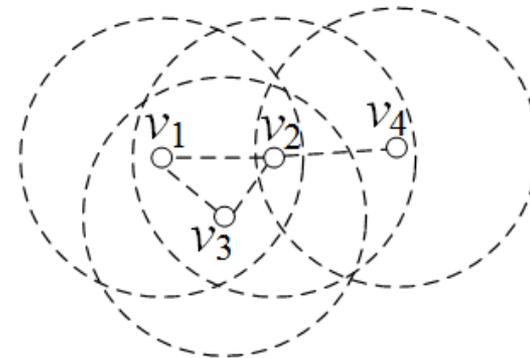


How to schedule a packet transmission?

➤ Collision free for packet #5

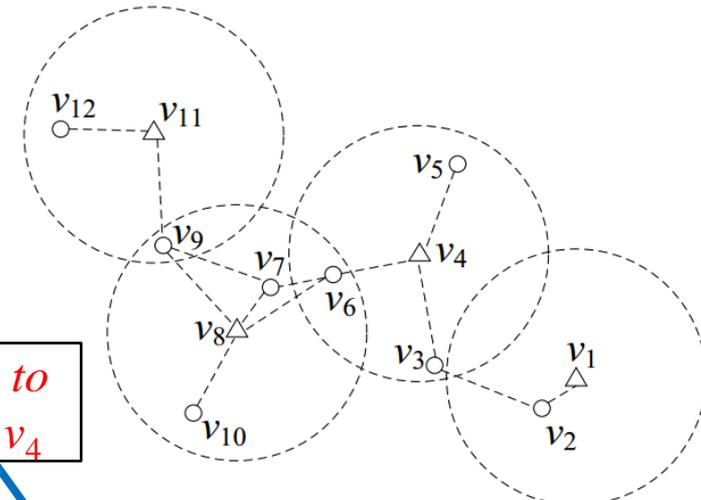


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- \square Newly scheduled TX slot



Scheduling in DOS

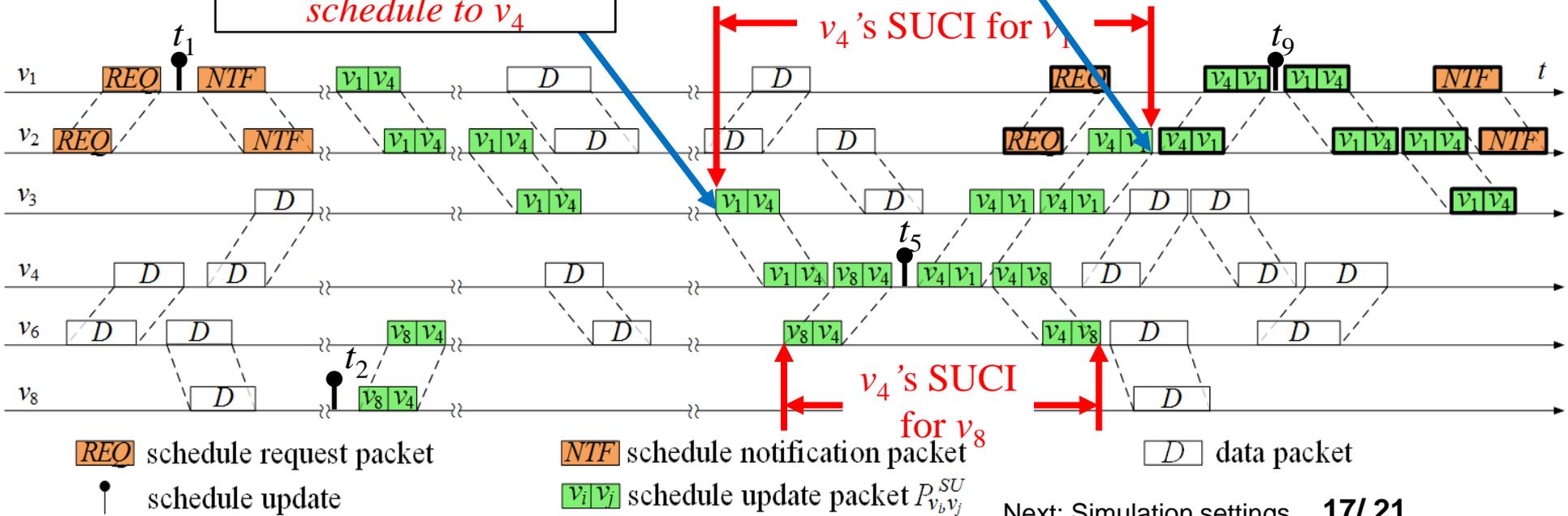
- Schedule update time
- Schedule update conflict interval (SUCI)
- Schedule notification packet
- Request packet
- Notification packet



○ cluster member △ cluster head

Deadline for v_1 to send schedule to v_4

Earliest time for v_1 to recv schedule from v_4

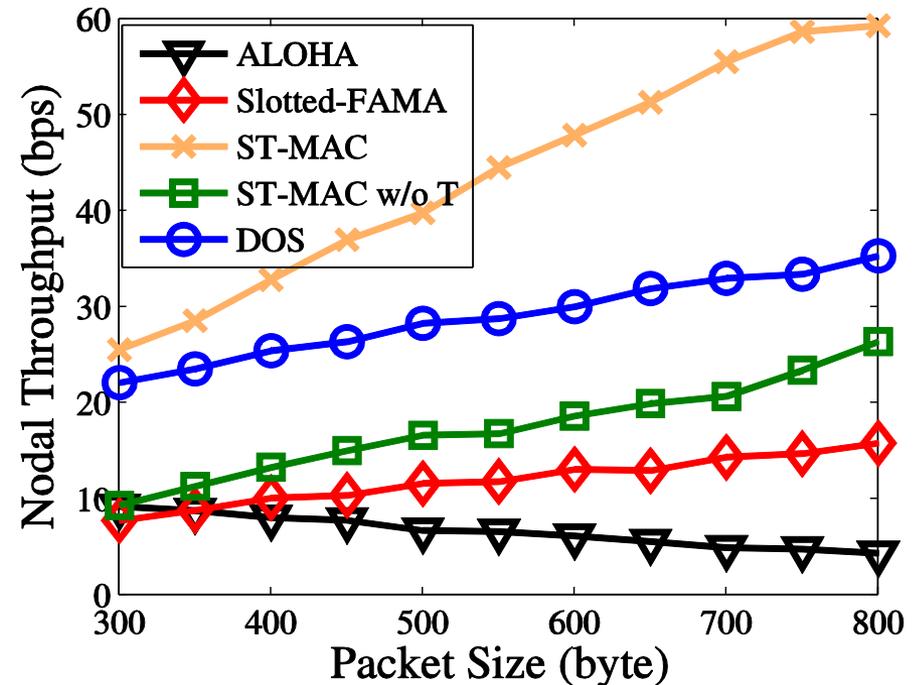
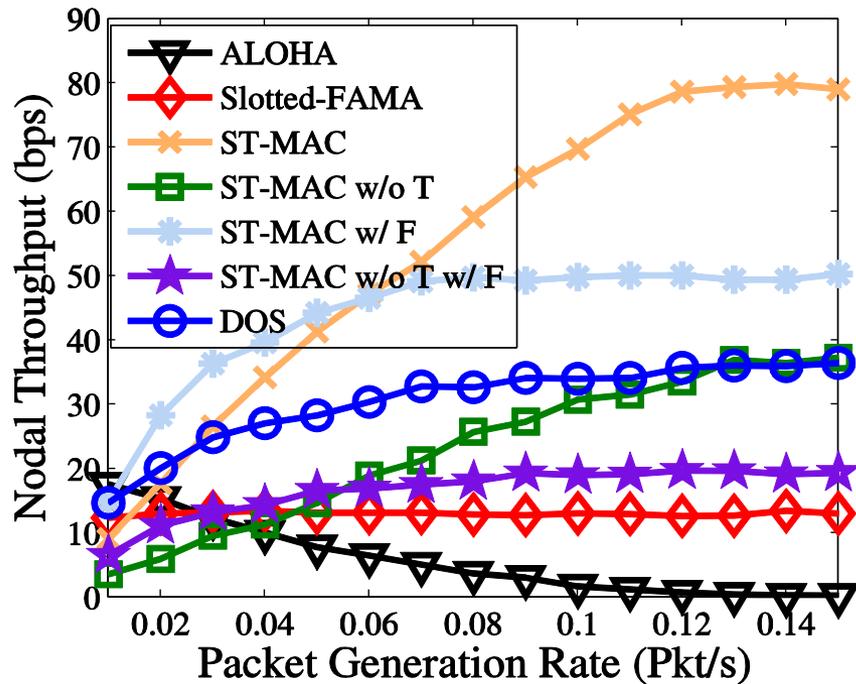


Simulation results

➤ Simulation settings

- Platform: Aqua-Sim
- Protocol comparison
 - ST-MAC and its three variants
 - Slotted-ALOHA
 - Slotted-FAMA
- Poisson traffic pattern
- Maximum transmission range 1100m
- Teledyne Benthos modem:
 - Transmission rate: 667 bps;
 - Preamble:1.5 s
- Network deployment: 80 nodes randomly distributed in $6000\text{m} \times 6000\text{m} \times 20\text{m}$ area; 80 sender/receiver pairs

Simulation results



DOS achieves comparable throughput to ST-MAC, a centralized solution, far outperforms others.

Summary

- Propose a distributed on-demand scheduling MAC
 - Collision-free: both data and **control** packets
 - Distributed: schedules locally, no global information required
 - On-demand: adapt to nodes' dynamic traffic
 - Pure-scheduling: No special requirement on modem device: CDMA, multi-channel, multi-user
 - Achieves high performance

- Future Work
 - Integrate DOS with a clustering algorithm capable to manage dynamic topology
 - Explore a cross-layer design approach to couple DOS with cluster-based routing

Thanks & Questions?